

Shared-Memory Parallelism and OpenMP

NAS Webinar

August 22, 2012

NASA Advanced Supercomputing Division

Outline



- Shared-memory parallelism
- What is OpenMP?
- OpenMP major components
 - Fork-join execution model
 - Worksharing constructs
 - Data sharing
 - Synchronization primitives
- Use of OpenMP
- Hybrid MPI + OpenMP model
- OpenMP tasking
- Performance considerations
- Future OpenMP extensions

Topics for the next webinar

Shared-Memory Parallelism



- Many modern parallel computers
 - A cluster of shared-memory nodes with multicore CPUs
- Size of shared-memory nodes getting larger
 - Increased number of cores (4, 8, 16 ...)
 - Many cores in new types of systems (such as GPUs, Intel MIC)
- Shared-memory programming
 - Access to the same, globally shared, address space
 - No need for explicit data communication
 - Possibility for maintaining sequential equivalency

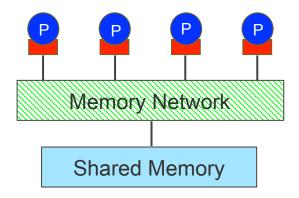


NASA's Pleiades Supercomputer

Shared-Memory Architecture

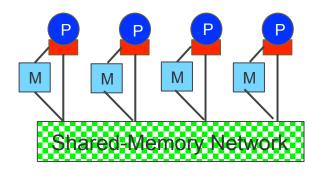


 Multiple processing units accessing global shared memory using a single address space



UMA: Uniform Memory Access

- "equidistant" access from all processors



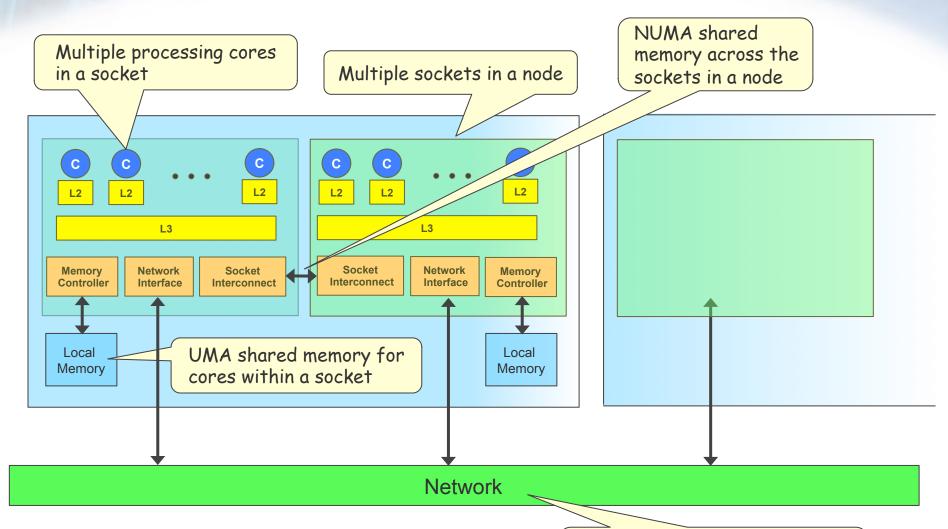
NUMA: Non-Uniform Memory Access

local memory versus remote memory

- Shared-memory systems easier to program
 - User responsible for synchronization of processors for correct data access and modification
- Scaling to large number of processors can be an issue

A Typical Supercomputer





Within a node: cache coherent, mix of UMA and NUMA shared memory

Distributed memory cluster of multi-socket nodes

Programming Approaches



- Thread-based approaches
 - Posix threads (low level)
 - OpenMP (de-facto standard)
 - Intel Thread Building Block
- Task-based approaches
 - Intel Cilk++
 - OpenMP 3.0
 - Grand Central Dispatch from Apple
- Others
 - Global arrays
 - Compiler auto-parallelization

What is OpenMP?



- A standard API to support shared-memory multiprocessing programming
 - Compiler directives and library routines for C/C++ and Fortran
 - Specification defined and maintained by the OpenMP Architecture Review Board
 - OpenMP 1.0 released in October 1997 for Fortran, 1998 for C/C++
 - Latest 3.1 released in July 2011
 - Implemented and supported by many compiler vendors
 - (Intel, PGI, IBM, Oracle, GCC, etc.)



Compiler Directives



- Special #pragma in C/C++, special comments in Fortran
- Often only enabled by a special compiler flag
- Program may be run sequentially when directives are ignored
- Examples of compiler directives

C/C++

```
#pragma omp parallel for for (i = 0; i < n; i++)
    a[i] = b[i] + c[i];
```

Fortran

```
!$omp parallel do
  do i = 1, n
    a(i) = b(i) + c(i)
  end do
!$omp end parallel do
```

Compile with "-openmp" (Intel compilers) to enable the compiler directives, otherwise they are treated as comments and the loop is run sequentially.

Advantages of OpenMP

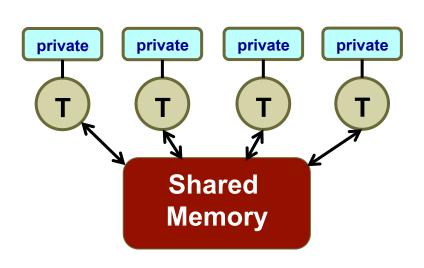


- Directive-based approach
 - Possible to write sequentially consistent code
 - Easier maintenance
- Global view of application memory space
 - Relatively faster program development
- Incremental parallelization
 - Piecemeal code development
 - Easier to program and debug
- When mixed with MPI
 - Maps well with multicore hybrid architectures

Major Components



- OpenMP thread
 - Execution entity with a stack and its private memory
 - Dynamically created and managed by the OpenMP runtime library
 - Access to shared memory
- Language components
 - Fork-join model for structured programming
 - Worksharing constructs for work distribution
 - Data sharing attributes
 - Synchronization primitives
- Runtime library routines
- Environment variables

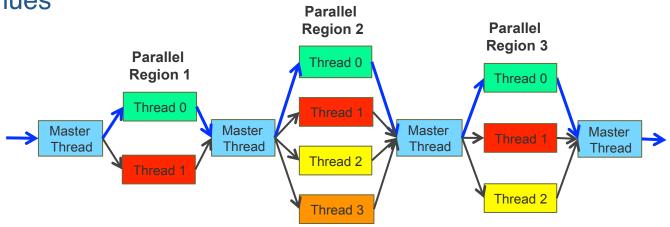


Execution Model



- Fork-join model
 - Program starts with a single (master) thread
 - Multiple threads are forked by the *master* thread at a **parallel** construct
 - The *master* thread is part of the new team of threads
 - Threads perform work in the parallel region
 - Worksharing constructs distribute work among threads
 - Threads may be synchronized with synchronization constructs

Threads join at the end of the **parallel** region and the *master* thread continues



Parallel Construct



- The fundamental construct to start parallel execution
 - Invocation of a team of threads
 - Code executed redundantly by every thread until a worksharing construct is encountered
 - Number of threads controlled via
 - The OMP_NUM_THREADS environment variable
 - A call to omp_set_num_threads(), or
 - The num_threads clause

```
omp_set_num_threads(4);
#pragma omp parallel private(myid)
{
   myid = omp_get_thread_num();
   printf("myid is %d\n", myid);
}
```

Worksharing Construct



- The construct to distribute work among threads
 - for (or do): used to split up loop iterations among the threads

```
#pragma omp for
for (i=0; i<n; i++) a[i] = b[i] + c[i];
```

- sections: assigning consecutive but independent code blocks to different threads (can be used to specify task parallelism)
 - Each code block is indicated by the section directive

```
#pragma omp sections
{ #pragma omp section
  work1();
  #pragma omp section
  work2();
}
```

Worksharing Construct (cont.)



single: specifying a code block executed by only one thread

```
#pragma omp single s = 0;
```

- There is an *implicit barrier* at the end of a worksharing construct
 - But can be suppressed with the "nowait" clause

```
#pragma omp for nowait
for (i=0; i<n; i++) a[i] = b[i] + c[i];
#pragma omp for
for (i=0; i<n; i++) d[i] = e[i] + f[i];</pre>
```

"nowait" suppresses the barrier at the end of the first **for** loop

- Master construct
 - code block executed by the master thread only, no barrier wait

Loop Scheduling



 Clause to define how loop iterations are distributed among threads of the team

```
#pragma omp for schedule(static)
for (i=0; i<n; i++) a[i] = b[i] + c[i];</pre>
```

- Loop scheduling kinds
 - **static**: for balanced workload, lowest overhead
 - Default for most compilers
 - **dynamic**: for unbalanced loop iterations
 - guided: for special monotonically increasing or decreasing workload
 - auto: compiler determines at runtime

Data Sharing



- Accessibility of variables by threads
 - **shared**: variable is shared by all threads in a team
 - private: variable is private to each thread
 - By default, variables are shared
 - With some exceptions, such as, loop variable is private
- Specifying data sharing attribute in a parallel region or worksharing construct
 - **shared** clause: for variables shared by threads
 - *private* clause: for variables private to each thread
 - reduction clause: combining private copies of a variable to the shared copy by an operator. Reduced final value is only guaranteed at a barrier.

Data Sharing (cont.)



An example

Threadprivate directive

- Special storage for global variables, private to each thread
- Specified at the variable declaration, valid throughout the program

```
static double c1, c2;
#pragma omp threadprivate(c1,c2)
```

Synchronization



- Barrier: wait until all of threads of a team have reached this point before continuing
 - Barrier construct specifies an explicit barrier
 - A worksharing construct has an implicit barrier at the end
- Critical construct
 - Code block executed by only one thread at a time, e.g., allows multiple threads to update shared data

```
#pragma omp critical updates the shared updates the shared variable "s" at a time

#pragma omp barrier

printf("sum = %g\n", s);

Ensure one thread updates the shared variable "s" at a time

All threads have done the update before the result is printed
```

Synchronization (cont.)



Atomic construct

- Update a shared variable atomically, can be more efficient than the critical construct if there is hardware support
- Only valid for scalar variable and a limited set of operations (+,*,-,...)

```
#pragma omp atomic
s = s + s_local;
#pragma omp barrier
printf("sum = %g\n", s);
```

- Other forms are also available:
 - "atomic read", "atomic write", "atomic capture"
- Locks
 - Similarly to critical but provided by the library routines and more flexible

Code Example: Sum of Squares



```
Needed for
                                               Forks off the threads and starts the
                             runtime routines
#include <omp.h>
                                               parallel execution; declares thread_id
long int sum = 0, loc_sum;
                                               and loc sum private
int thread_id;
#pragma omp parallel private(thread_id, loc_sum)
                                                       Each thread
                                                       retrieves its own id
  loc_sum = 0;
  thread_id = omp_get_thread_num();
                                                      Worksharing construct
  #pragma omp for schedule(static)
                                                      distributes the work
  for(i = 0; i < N; i++)
                                                       Each thread prints its
     loc_sum = loc_sum + i * i;
                                                       id and local sum
  printf("Thread %d: loc_sum = %ld\n", thread_id, loc_sum);
  #pragma omp critical
                                              Threads cooperate to update
  sum = sum + loc sum;
                                             the shared variable one by one
printf("sum = %ld\n",sum);
                                                  Master thread prints result
```

Use of OpenMP on Pleiades



Basic steps

- Select a compiler:

```
module load comp-intel/11.1.072
```

Compile codes with flags that enable OpenMP

```
icc -o s1.x -O3 -openmp squares.c
```

- Set the number of threads to be used

```
setenv OMP_NUM_THREADS 8
```

- Run the executable

```
./s1.x
```

For details see

http://www.nas.nasa.gov/hecc/support/kb/With-OpenMP_103.html

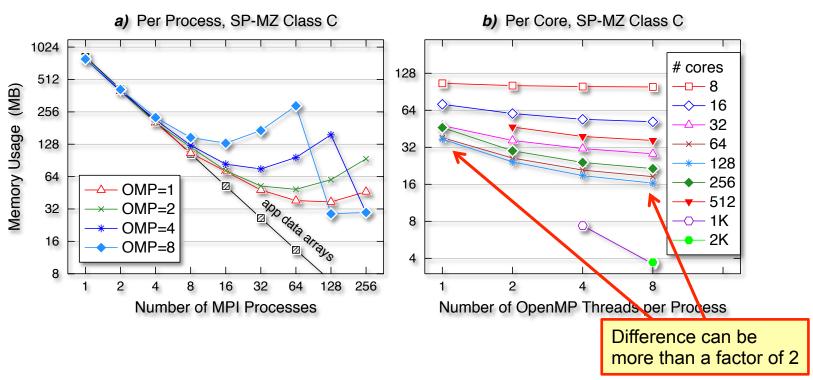
Hybrid MPI + OpenMP



- The hybrid model
 - OpenMP works in the memory space of each MPI process
 - Shared memory within each MPI process but distributed memory across MPI processes
- Advantages of the hybrid model
 - Maps well to many hardware architectures, including Pleiades
 - MPI for communication between distributed-memory nodes
 - OpenMP for shared-memory parallelism with a node
 - Can achieve good scalability when not possible with pure MPI
 - A hybrid code may consume less memory than a pure MPI code

Memory Usage of Hybrid Codes





The SP-MZ benchmark on the SGI Altix ICE

Hybrid Programming



Approaches

- Common approach
 - MPI for parallelism at coarser level, OpenMP at finer level
 - No MPI calls inside OpenMP parallel regions
- Mixed approach
 - MPI routines may be called inside OpenMP parallel regions
 - Requires the MPI library to be thread-safe (MPI_THREAD_MULTIPLE)
- Program development
 - MPI and OpenMP can be developed separately

Hybrid Code Example: Computing Pi



```
include "mpif.h"
                                                   Get rank and size of
integer myid, numprocs, ierr, n, i
                                                   MPI processes
real(8) h, sum, mypi, pi
call MPI_Comm_rank(MPI_COMM_WORLD,myid,ierr)
call MPI_Comm_size(MPI_COMM_WORLD,numprocs,ierr)
n = 100000
                                                  Use OpenMP to
h = 1.0d0/n
                                                  compute partial sum
sum = 0.0d0
!$omp parallel do private(x) reduction(+:sum)
do i = myid+1, n, numprocs
 x = h * (i - 0.5d0)
  sum = sum + 4.0d0 / (1.0d0 + x * x)
                                                  Reduce the final result
                                                  from all MPI processes
end do
mypi = h * sum
call MPI_Reduce(mypi,pi,1,MPI_REAL8,MPI_SUM, &
                 O,MPI_COMM_WORLD,ierr)
if (myid.eq.0) print *,"pi is ",pi
                                                 Rank 0 prints result
```

Use of MPI + OpenMP on Pleiades



Basic steps

- Select a compiler and an MPI library: module load comp-intel/11.1.072 mpi-sqi/mpt.2.06r6
- Compile codes with flags that enable OpenMP and link with MPI library ifort -o s2.x -O3 -openmp pi_hybrid.f90 -lmpi
- Set thread and process binding flags (for performance reason)
 setenv MPI_DSM_DISTRIBUTE
 setenv MPI_OPENMP_INTEROP
- Set the number of threads to be used setenv OMP_NUM_THREADS 4
- Run the executable (with 2 MPI processes, 4 OpenMP threads/process)
 mpiexec -np 2 ./s2.x
- For details see

http://www.nas.nasa.gov/hecc/support/kb/52/

OpenMP Performance Issues



- Why is my OpenMP code not scaling? Possible issues:
 - Overhead of OpenMP constructs
 - Granularity of work units
 - Remote data access and NUMA effect
 - Load imbalance
 - False sharing of cache
 - Poor resource utilization
- We will discuss these issues and possible solutions together with other advanced OpenMP topics in the next webinar

References



- OpenMP specifications
 - www.openmp.org/wp/openmp-specifications/
- Resources
 - www.openmp.org/wp/resources/
 - www.compunity.org/
- Benchmarks
 - OpenMP Microbenchmarks from EPCC (www.epcc.ed.ac.uk/research/openmpbench)
 - NAS Parallel Benchmarks (www.nas.nasa.gov/publications/npb.html)
- Porting applications to Pleiades
 - www.nas.nasa.gov/hecc/support/kb/52/
 - www.nas.nasa.gov/hecc/support/kb/60/